1-Bit Matrix Completion

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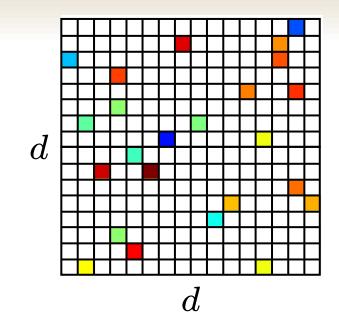
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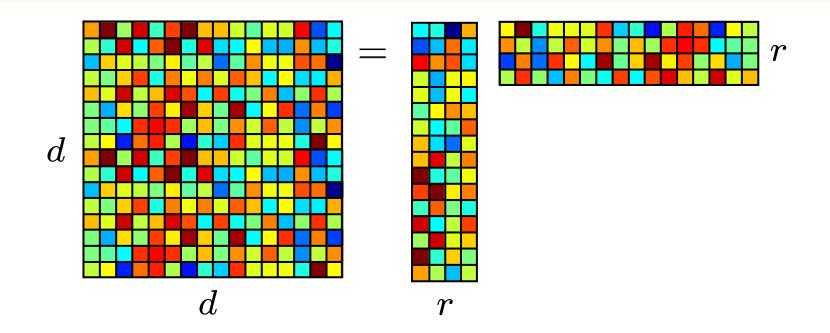


Matrix Completion



Under what assumptions can we recover the original matrix?

Low-Rank Matrices



Singular value decomposition:

$$M = U\Sigma V^*$$

 $\approx dr \ll d^2$

degrees of freedom

Low-Rank Matrix Recovery

Given:

- a $d \times d$ matrix M of rank r
- samples of $M \, {\rm on}$ the set $\ : \ Y = M$

How can we recover M?

$$\widehat{M} = \operatorname*{arg\,inf}_{X:X = Y} \operatorname{rank}(X)$$

Can we replace this with something computationally feasible?

Nuclear Norm Minimization

Convex relaxation!

Replace rank(X) with
$$||X||_* = \sum_{j=1}^d |\sigma_j|$$

$$\widehat{M} = \underset{X:X = Y}{\operatorname{arg inf}} \|X\|_*$$

If $| = O(r d \log d)$, this procedure can recover M!

Applications

- Collaborative Filtering (aka the "Netflix Problem")
- Recovery of incomplete survey data
- Analysis of voting data
- Sensor localization
- Quantum state tomography

Matrix Completion in Practice

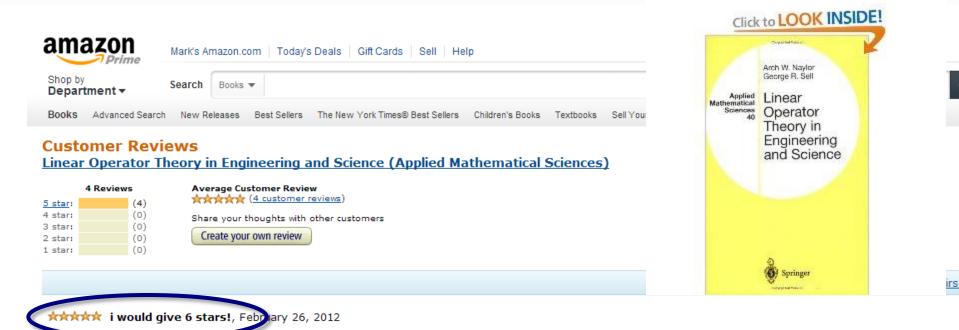
Noise

$$Y = (M + Z)$$

• Quantization

- Netflix/Amazon: Ratings are integers between 1 and 5
- Survey responses: True/False, Yes/No, Agree/Disagree
- Voting data: Yea/Nay
- Quantum state tomography: Binary outcomes

What's the Problem?



Amazon Verified Purchase (What's this?)

By Chen

This review is from: Linear Operator Theory in Engineering and Science (Applied Mathematical Sciences) (Paperback)

I'm doing a PhD in econometrics and I need to apply operator theories in constructing a linear or nonlinear operator to help explain individual economic behaviour. This book contains numerous useful ideas and applications with exercises thoroughly designed; one of the questions in the exercise gave me an idea of creating a matrix for describing a nonlinear operator. That question asks for a matrix that describes a second order differential operator and that gave me an idea that taylor series approximation can be used to linearise a nonlinear operator and hence a nonlinear operator may also be described by a matrix.

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1-Bit Matrix Completion

Extreme case

$$Y = \operatorname{sign}(M)$$

Claim: Recovering M from Y is impossible!

No matter how many samples we obtain, all we can learn is whether $\lambda>0\,$ or $\,\lambda<0\,$

Is There Any Hope?

If we consider a noisy version of the problem, recovery becomes feasible!

$$Y = \operatorname{sign}(M + Z)$$

$$M + Z = \begin{bmatrix} \lambda + Z_{1,1} & \lambda + Z_{1,2} & \lambda + Z_{1,3} & \lambda + Z_{1,4} \\ \lambda + Z_{2,1} & \lambda + Z_{2,2} & \lambda + Z_{2,3} & \lambda + Z_{2,4} \\ \lambda + Z_{3,1} & \lambda + Z_{3,2} & \lambda + Z_{3,3} & \lambda + Z_{3,4} \\ \lambda + Z_{4,1} & \lambda + Z_{4,2} & \lambda + Z_{4,3} & \lambda + Z_{4,4} \end{bmatrix}$$

Fraction of positive/negative observations tells us something about λ

Example of the power of *dithering*

Observation Model

For $(i, j) \in$ we observe

$$Y_{i,j} = \begin{cases} +1 & \text{with probability } f(M_{i,j}) \\ -1 & \text{with probability } 1 - f(M_{i,j}) \end{cases}$$

If f behaves like a CDF, then this is equivalent to

$$Y_{i,j} = \operatorname{sign}(M_{i,j} + Z_{i,j})$$

where $Z_{i,j}$ is drawn according to a suitable distribution

We will assume that is drawn uniformly at random

Examples

• Logistic regression / Logistic noise

$$f(x) = rac{e^x}{1 + e^x}$$

 $Z_{i,j} \sim ext{logistic distribution}$

• Probit regression / Gaussian noise

$$f(x) = \Phi(x/\sigma)$$

$$Z_{i,j} \sim \mathcal{N}(0,\sigma^2)$$

Maximum Likelihood Estimation

Log-likelihood function:

$$F(X) = \sum_{(i,j)\in +} \log(f(X_{i,j})) + \sum_{(i,j)\in -} \log(1 - f(X_{i,j}))$$

$$\widehat{M} = \operatorname*{arg\,max}_{X} F(X)$$

s.t.
$$\frac{1}{d\alpha} \|X\|_{*} \leq \sqrt{r}$$
$$\|X\|_{\infty} \leq \alpha$$

Recovery of the Matrix

Theorem (Upper bound achieved by convex ML estimator) Assume that $\frac{1}{d\alpha} ||M||_* \leq \sqrt{r}$ and $||M||_{\infty} \leq \alpha$. If is chosen at random with $\mathbb{E}|$ $|=m > d \log d$, then with high probability

$$\frac{1}{d^2} \|\widehat{M} - M\|_F^2 \le C\alpha L_\alpha \beta_\alpha \sqrt{\frac{rd}{m}}$$

where

$$L_{\alpha} := \sup_{|x| \le \alpha} \frac{|f'(x)|}{f(x)(1 - f(x))} \qquad \beta_{\alpha} := \sup_{|x| \le \alpha} \frac{f(x)(1 - f(x))}{(f'(x))^2}$$

Is this bound tight?

Recovery of the Matrix

Theorem (Upper bound achieved by convex ML estimator) Assume that $\frac{1}{d\alpha} ||M||_* \leq \sqrt{r}$ and $||M||_{\infty} \leq \alpha$. If is chosen at random with $\mathbb{E}||=m > d \log d$, then with high probability

$$\frac{1}{d^2} \|\widehat{M} - M\|_F^2 \le C\alpha L_\alpha \beta_\alpha \sqrt{\frac{rd}{m}}$$

Theorem (Lower bound on any estimator)

For any recovery algorithm \widehat{M} there exist M satisfying the assumptions above such that for any set with | = m, we have (under mild technical assumptions) that

$$\mathbb{E}\left[\frac{1}{d^2}\|\widehat{M} - M\|_F^2\right] \ge c\alpha \sqrt{\beta_{\frac{3}{4}\alpha}} \sqrt{\frac{rd}{m}}$$

Tiny Sketch of Proof of Upper Bound

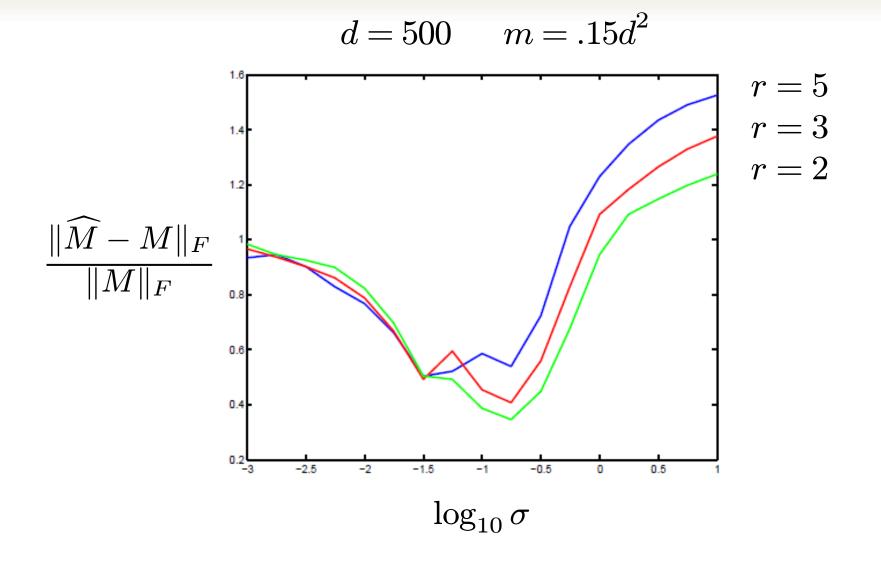
Recall that we maximize the log-likelihood F(X)

- For a fixed matrix X, $\mathbb{E}\left[F(M) F(X)\right] = c \cdot D(f(X)||f(M))$
- Lemma: Let $K = \{X : \frac{1}{d\alpha} ||X||_* \le \sqrt{r}\}$. With high probability, $\sup_{X \in K} |F(X) \mathbb{E}F(X)| \le \delta$
- By definition, $F(\widehat{M}) \ge F(M)$

$$0 \ge F(M) - F(\widehat{M})$$
$$\ge \mathbb{E}\left[F(M) - F(\widehat{M})\right] - 2\delta$$
$$= c \cdot D(f(\widehat{M})||f(M)) - 2\delta$$

• Thus, $D(f(\widehat{M})||f(M)) \leq \frac{2}{c}\delta$

Synthetic Simulations



MovieLens Data Set

- 100,000 movie ratings on a scale from 1 to 5
- Convert to binary outcomes by comparing each rating to the average rating in the data set
- Evaluate by checking if we predict the correct sign
- Training on 95,000 ratings and testing on remainder
 - "standard" matrix completion: 60% accuracy
 - 1: 64%
 2: 56%
 3: 44%
 4: 65%
 5: 74%
 - 1-bit matrix completion: 73% accuracy
 - 1: 79%
 2: 73%
 3: 58%
 4: 75%
 5: 89%

Conclusions

- 1-bit matrix completion is hard!
- What did you really expect?
- Sometimes 1-bit is all we can get...
- We have algorithms that are near optimal
- Open questions
 - Are there simpler/better/faster/stronger algorithms?
 - What about 2.32-bit matrix completion?

Thank You!