

Compressive Measurements for Signal Acquisition and Processing

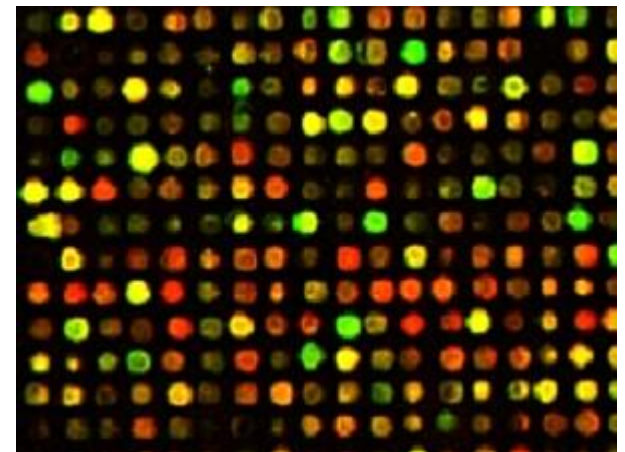
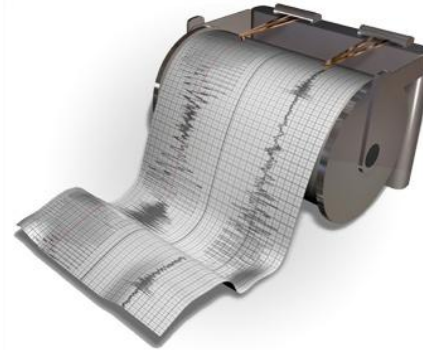
Mark Davenport



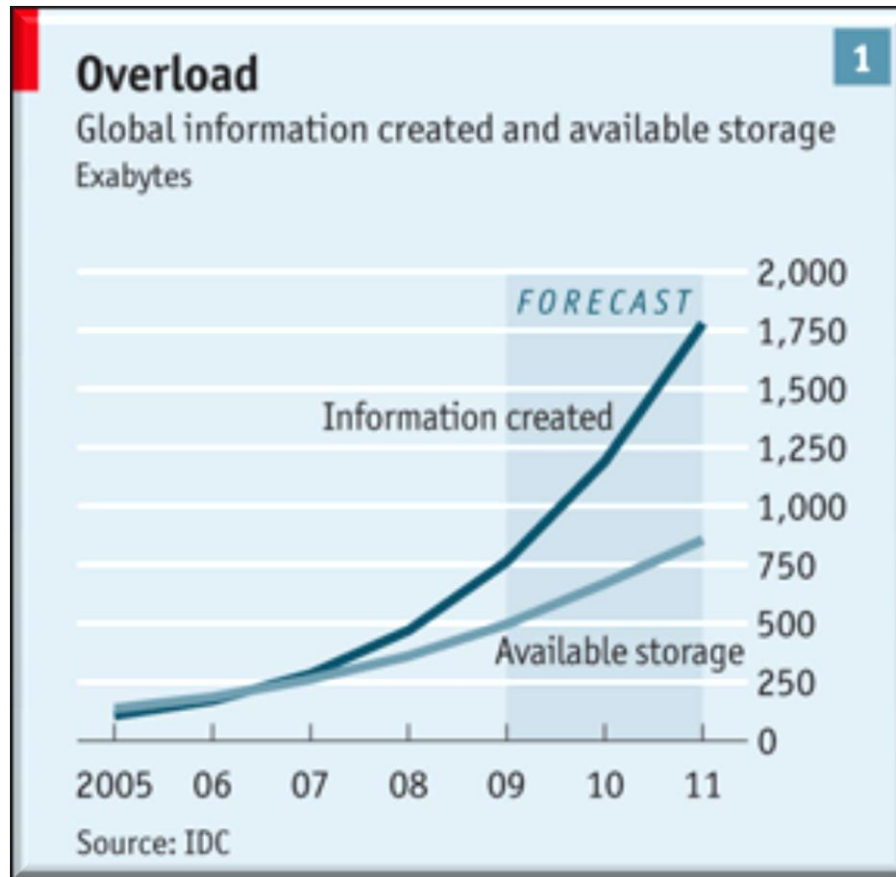
Rice University
ECE Department



Sensor Explosion



Data Deluge



By 2011, ½ of digital universe will have no home

Dimensionality Overload

~~How can we get as much data as possible into the digital domain?~~



How can we avoid having to acquire so much data?

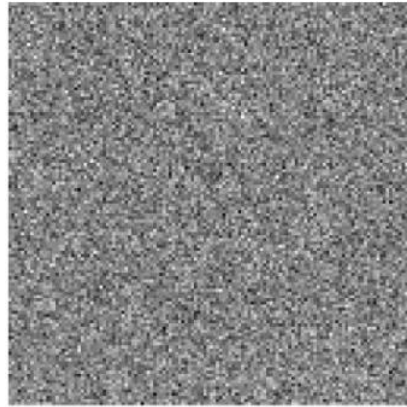
~~How can we extract as much information as possible from a limited amount of data?~~



How can we extract any information at all from a massive amount of high-dimensional data?

Dimensionality Reduction

Data is rarely intrinsically high-dimensional



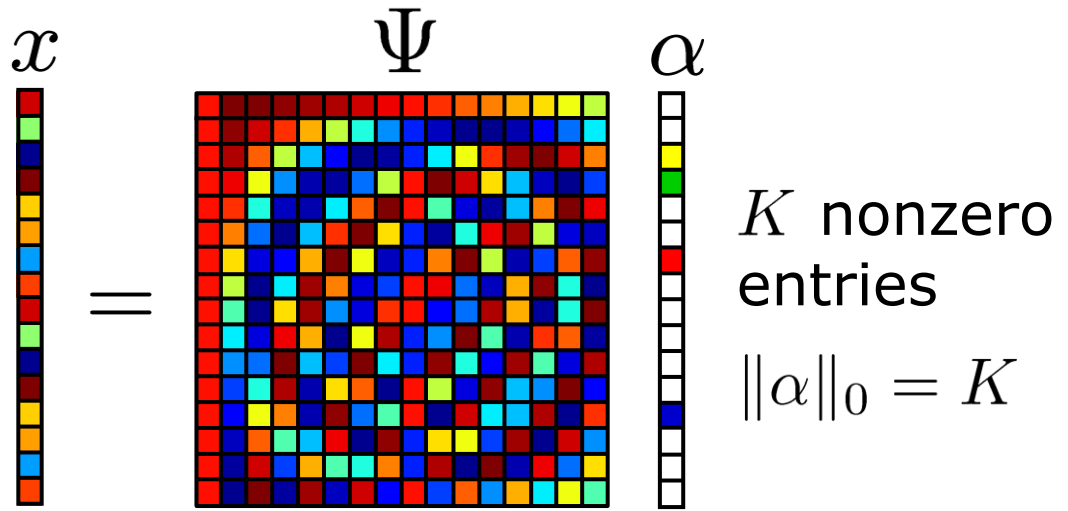
Signals often obey ***low-dimensional models***

- sparsity
- manifolds
- low-rank matrices

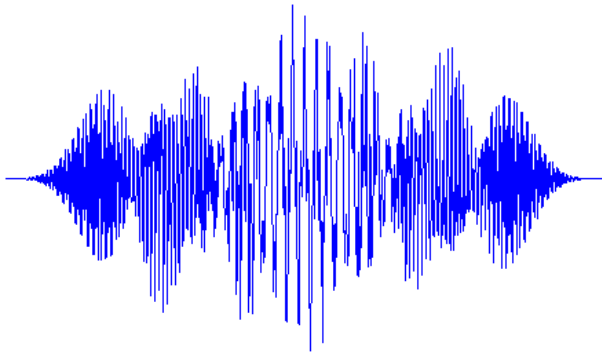
The intrinsic dimension K can be much less than the ambient dimension N , which enables ***dimensionality reduction***

Sparsity

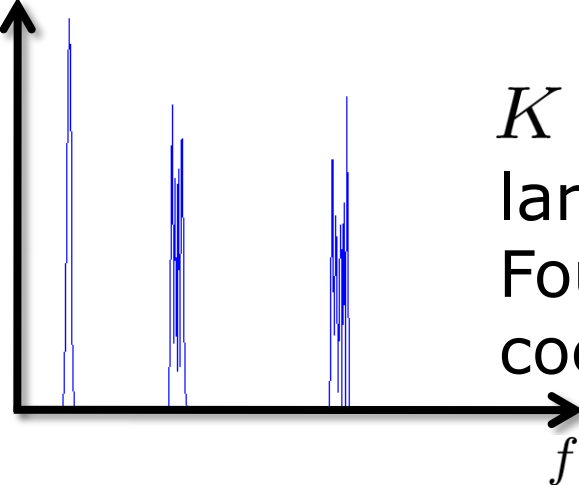
$$x = \sum_{j=1}^N \alpha_j \psi_j$$
$$= \Psi \alpha$$



N
samples



$X(f)$

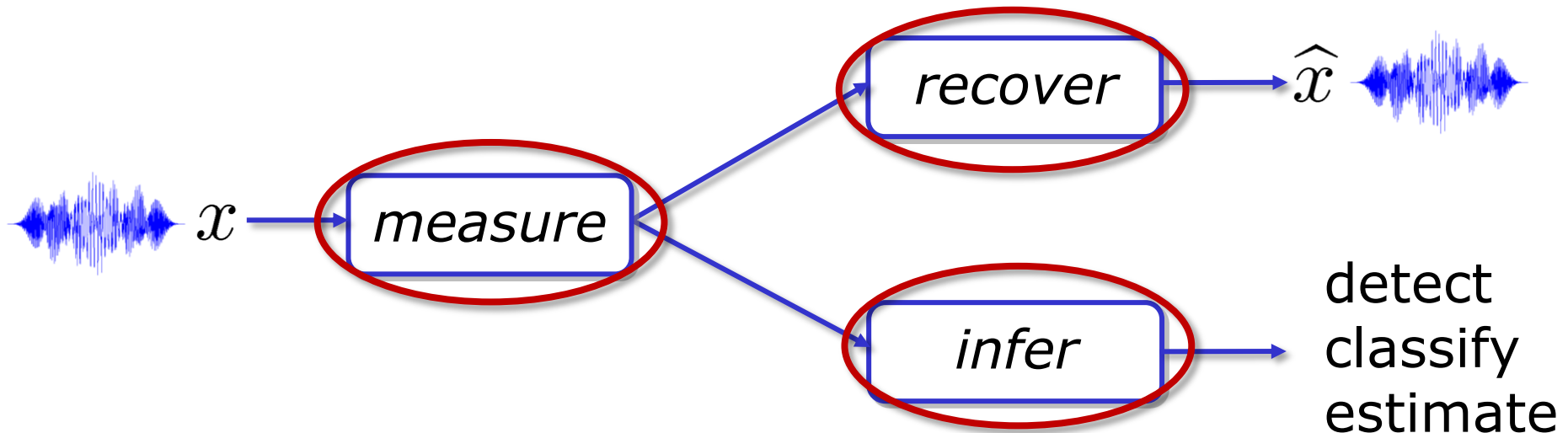


$K \ll N$
large
Fourier
coefficients

Compressive Signal Processing

How can we exploit low-dimensional models in the design of signal processing algorithms?

We would like to operate at the *intrinsic dimension* at all stages of the DSP pipeline



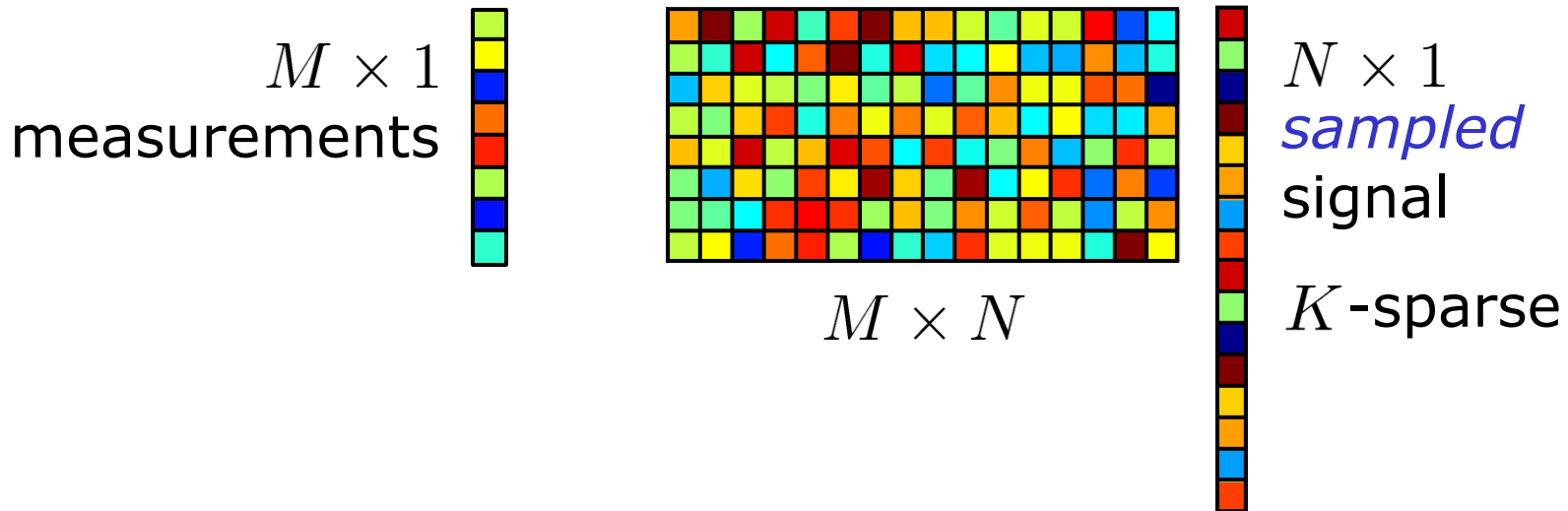
Compressive Measurements

Compressive Measurements

Compressive sensing [Donoho; Candes, Romberg, Tao – 2004]

Replace samples with general *linear measurements*

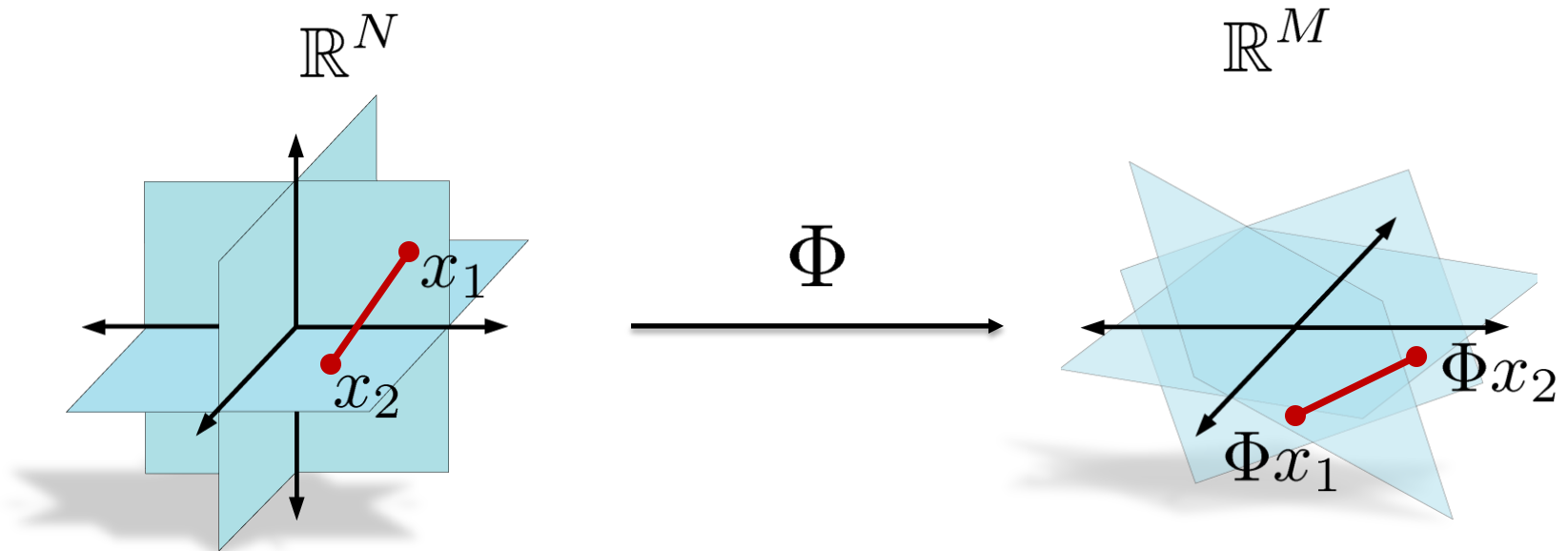
$$y = \Phi x$$



$$M = O(K \log(N/K))$$

Restricted Isometry Property (RIP)

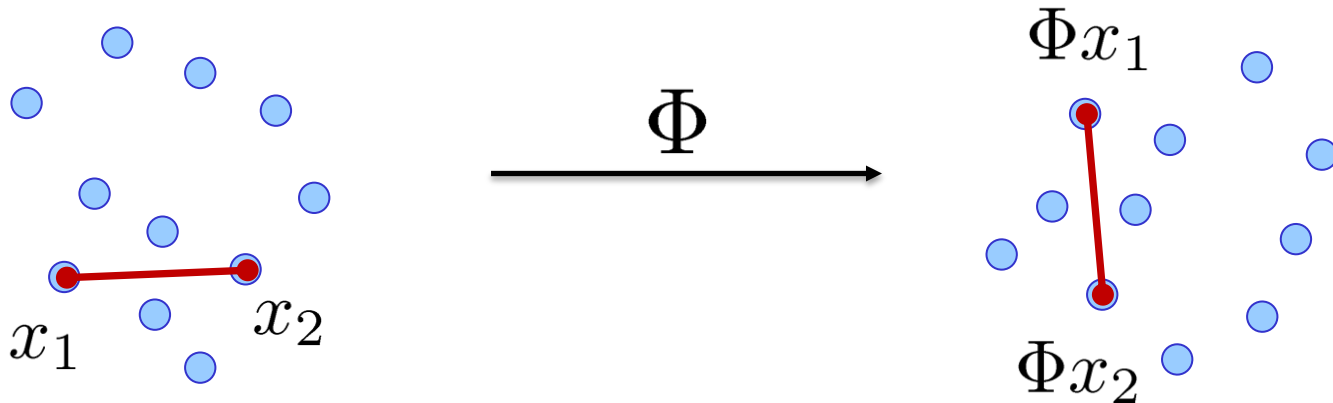
$$1 - \delta \leq \frac{\|\Phi x_1 - \Phi x_2\|_2^2}{\|x_1 - x_2\|_2^2} \leq 1 + \delta \quad \|x_1\|_0, \|x_2\|_0 \leq K$$



$$1 - \delta \leq \frac{\|\Phi x\|_2^2}{\|x\|_2^2} \leq 1 + \delta \quad \|x\|_0 \leq 2K$$

Johnson-Lindenstrauss Lemma

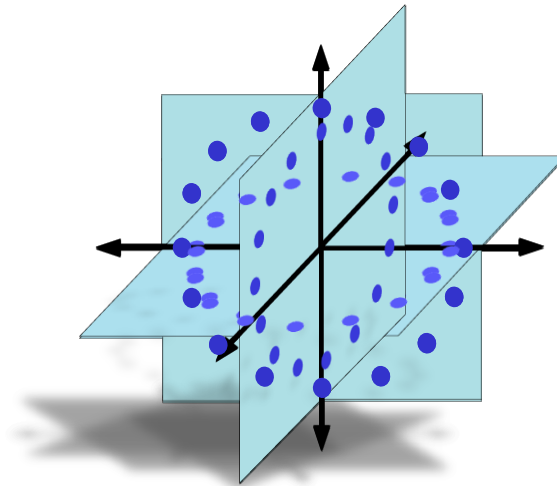
- Stable projection of a discrete set of P points



- Pick Φ at *random* using a *sub-Gaussian* distribution
- For any fixed x , $\|\Phi x\|_2$ concentrates around $\|x\|_2$ with (exponentially) high probability
- We preserve the length of all $O(P^2)$ difference vectors simultaneously if $M = O(\log P^2) = O(\log P)$.

JL Lemma Meets RIP

$$1 - \delta \leq \frac{\|\Phi x\|_2^2}{\|x\|_2^2} \leq 1 + \delta \quad \|x\|_0 \leq 2K$$



$$P = O\left(\left(\frac{N}{K}\right)^K\right) \longrightarrow M = O\left(K \log\left(\frac{N}{K}\right)\right)$$

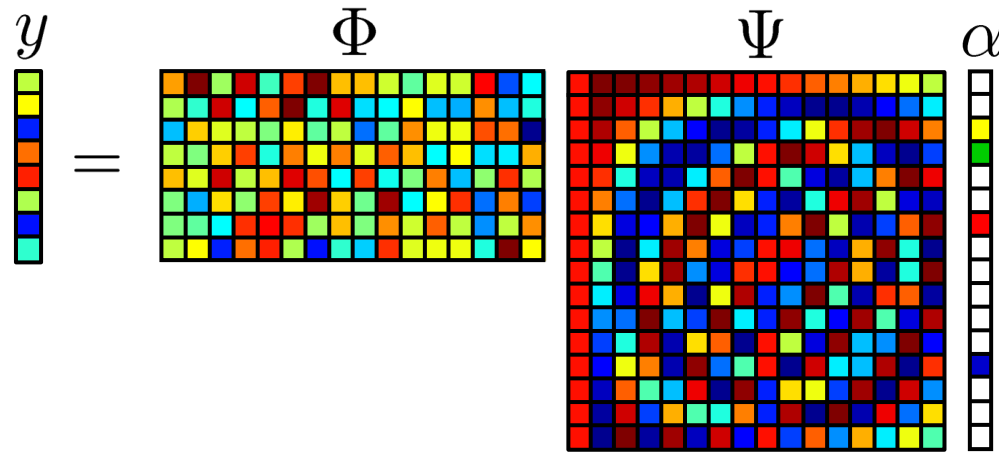
Hallmarks of Random Measurements

Stable

Φ will preserve information, be robust to noise

Universal

Φ will work with **any** fixed orthonormal basis (w.h.p.)

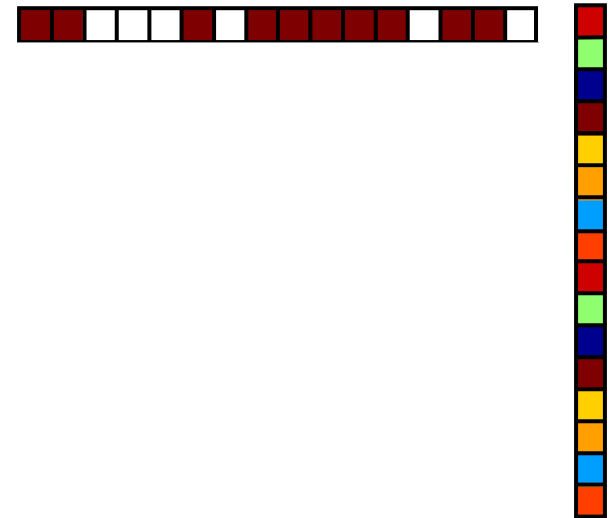
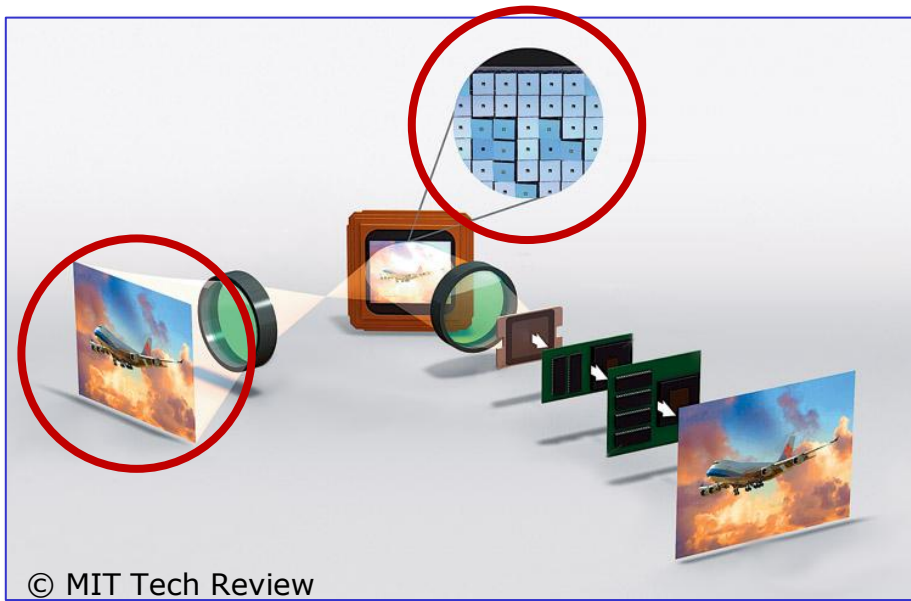


Democratic

Each measurement has "equal weight"

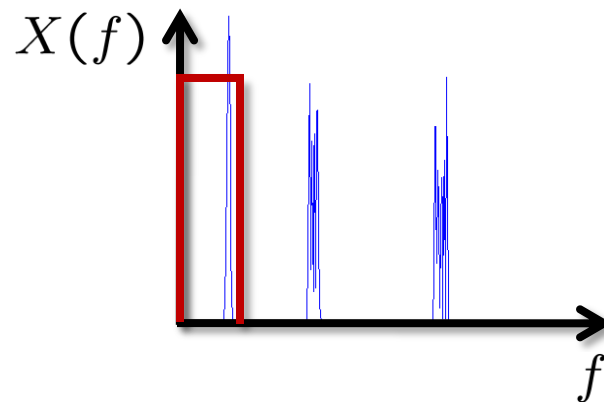
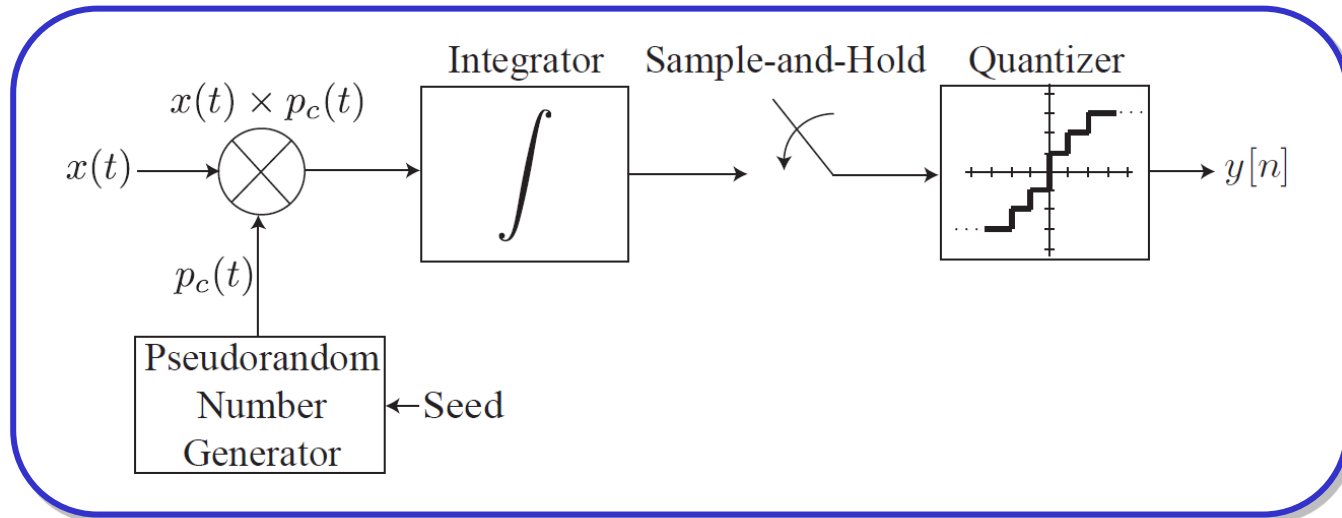
Compressive Measurements: Imaging

Rice "single-pixel camera"



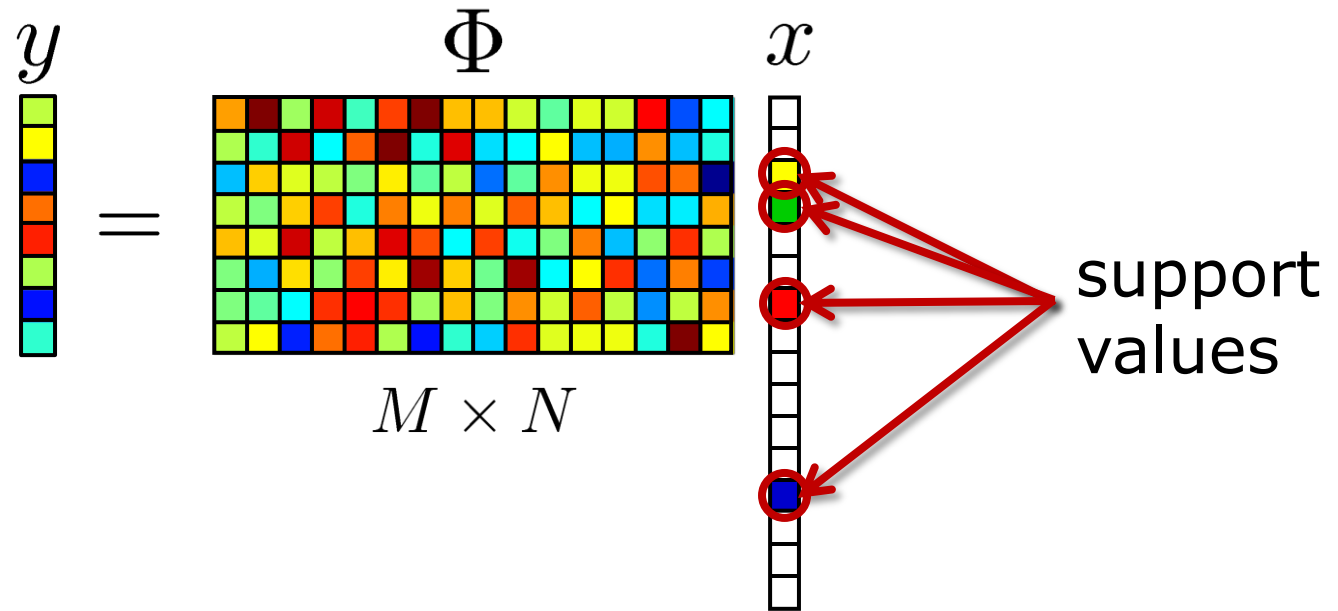
Compressive Measurements: ADCs

“Random demodulator”



Signal Acquisition and Recovery

Sparse Signal Recovery



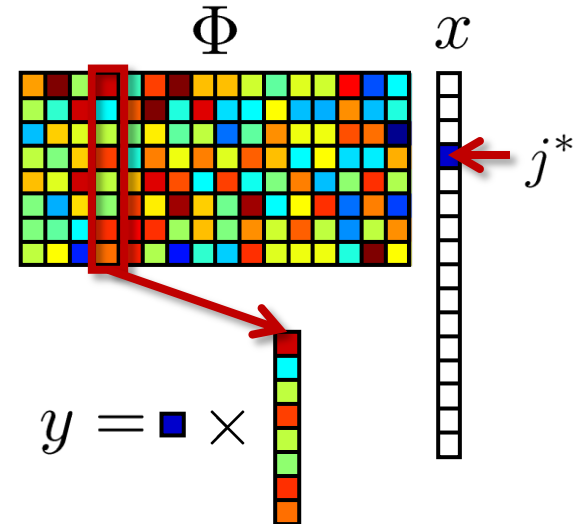
- Optimization / ℓ_1 -minimization
- Greedy algorithms
 - matching pursuit
 - orthogonal matching pursuit (OMP)
 - regularized OMP
 - CoSaMP, Subspace Pursuit, IHT, ...

Orthogonal Matching Pursuit

OMP selects one index at a time

Iteration 1:

$$j^* = \arg \max_j |\langle y, \Phi_j \rangle|$$



If Φ satisfies the RIP of order $\|u \pm v\|_0$, then

$$|\langle \Phi u, \Phi v \rangle - \langle u, v \rangle| \leq \delta \|u\|_2 \|v\|_2$$

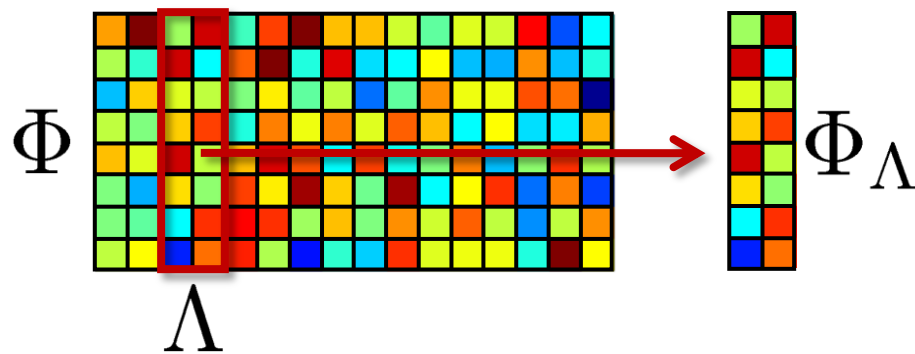
Set $u = x$ and $v = e_j$

$$|\langle y, \Phi_j \rangle - x_j| \leq \delta \|x\|_2$$

Orthogonal Matching Pursuit

Subsequent Iterations:

$$j^* = \arg \max_j |\langle Py, P\Phi_j \rangle|$$



$$P = I - \underbrace{\Phi_\Lambda \Phi_\Lambda^\dagger}_{\text{Projection onto } \mathcal{R}(\Phi_\Lambda)}$$

$$P\Phi_\Lambda = 0 \quad \longrightarrow \quad P\Phi x = P\Phi x_{\Lambda^c}$$

Interference Cancellation

Lemma

If Φ satisfies the RIP of order K , then

$$\left(1 - \frac{\delta}{1 - \delta}\right) \|x\|_2^2 \leq \|P\Phi x\|_2^2 \leq (1 + \delta) \|x\|_2^2$$

for all x such that $\|x\|_0 \leq K - |\Lambda|$ and $\text{supp}(x) \cap \Lambda = \emptyset$.

$$\rightarrow |\langle Py, P\Phi_j \rangle - x_j| \leq \frac{\delta}{1 - \delta} \|x_{\Lambda^c}\|_2$$

Orthogonal Matching Pursuit

Theorem

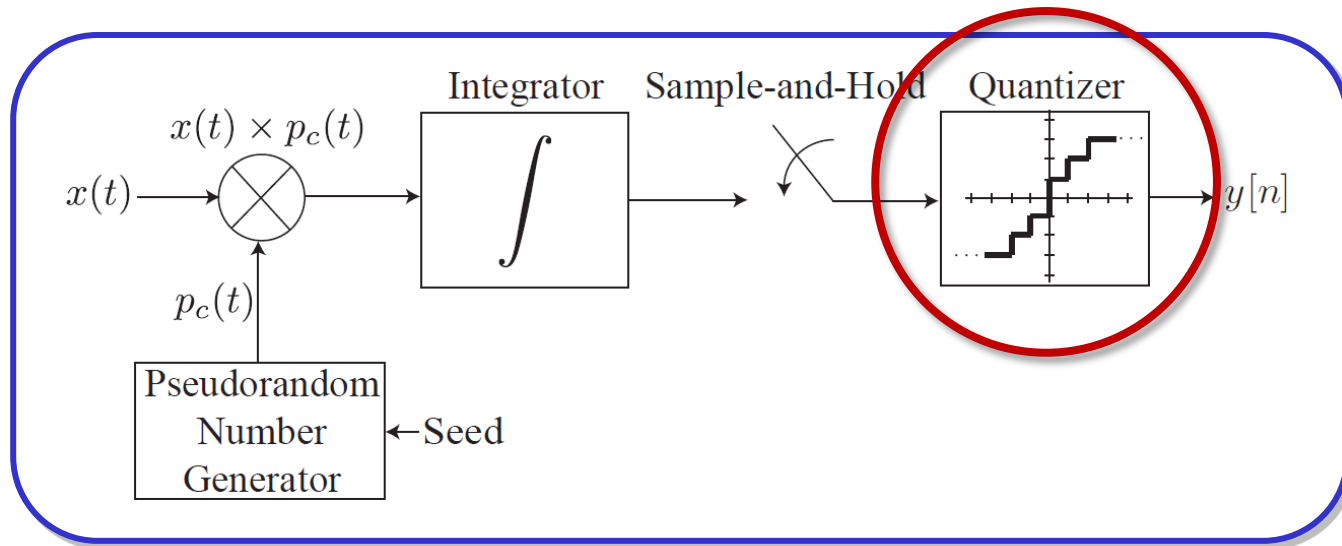
Suppose x is K -sparse and $y = \Phi x$.

If Φ satisfies the RIP of order $K + 1$ with constant $\delta < \frac{1}{3\sqrt{K}}$, then the j^* identified at each iteration will be a nonzero entry of x .

➡ Exact recovery after K iterations.

Argument provides simplified proofs for other orthogonal greedy algorithms (e.g. ROMP) that are robust to noise

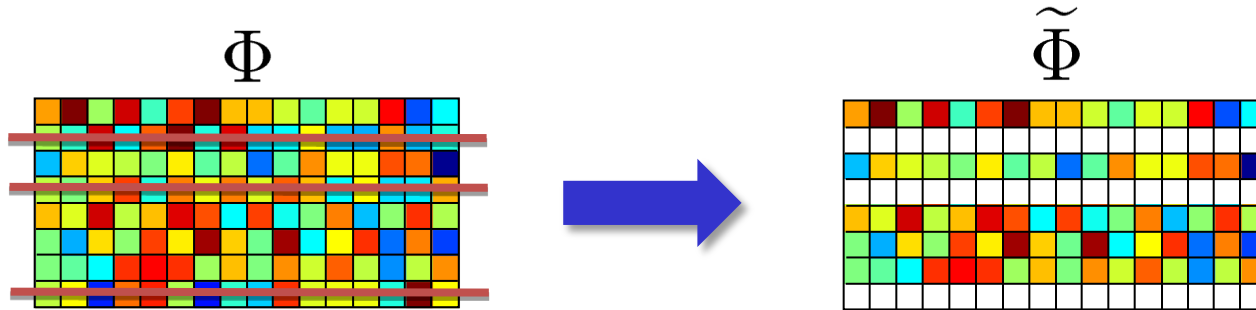
Signal Recovery with Quantization



- Most algorithms are designed for *bounded errors*
- Finite-range quantization leads to *saturation* and *unbounded errors*
- Being able to handle saturated measurements is *critical* in any real-world system

Saturation Strategies

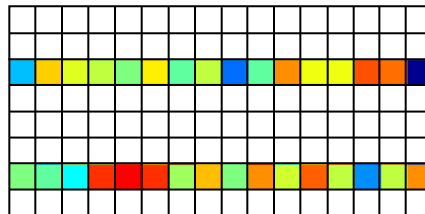
- **Rejection:** Ignore saturated measurements



- **Consistency:** Retain saturated measurements. Use them only as inequality constraints on the recovered signal
- If the rejection approach works, the consistency approach should automatically do better

Rejection and Democracy

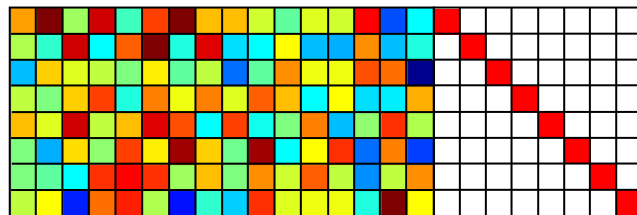
- The RIP is ***not sufficient*** for the rejection approach
- Example: $\Phi = I$
 - perfect isometry
 - *every* measurement must be kept
- We would like to be able to say that ***any*** submatrix of Φ with sufficiently many rows will still satisfy the RIP



- Strong, ***adversarial*** form of “democracy”

Sketch of Proof

- Step 1: Concatenate the identity to Φ



Theorem:

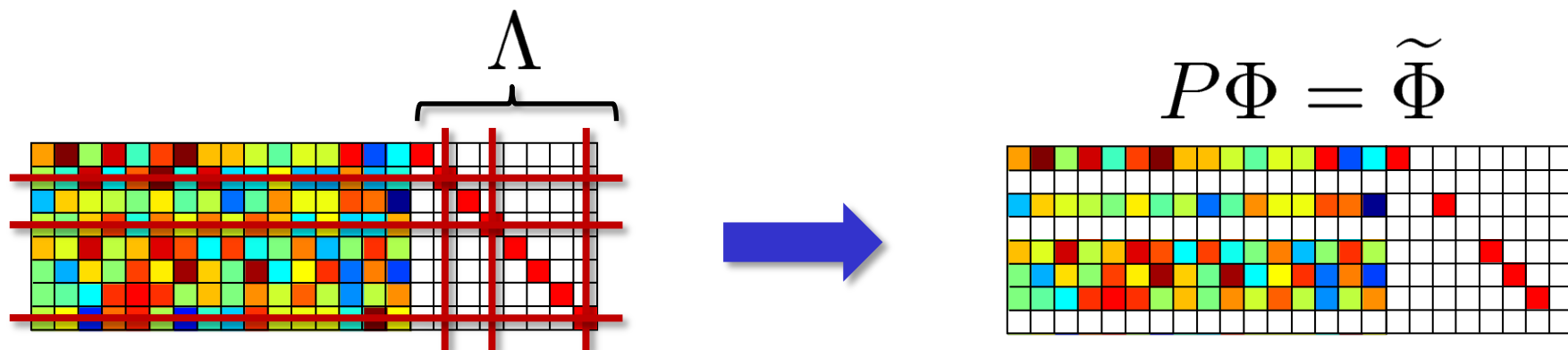
If Φ is a sub-Gaussian matrix with

$$M = O \left(K \log \left(\frac{N}{K} \right) \right)$$

then $[\Phi \ I]$ satisfies the RIP of order K with probability at least $1 - 3e^{-CM}$.

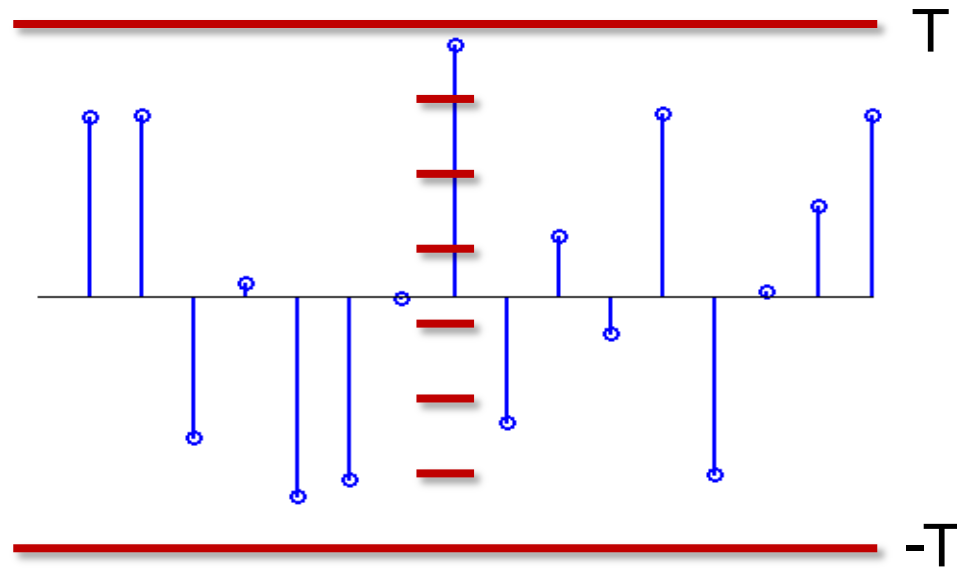
Sketch of Proof

- Step 2: Combine with the “interference cancellation” lemma



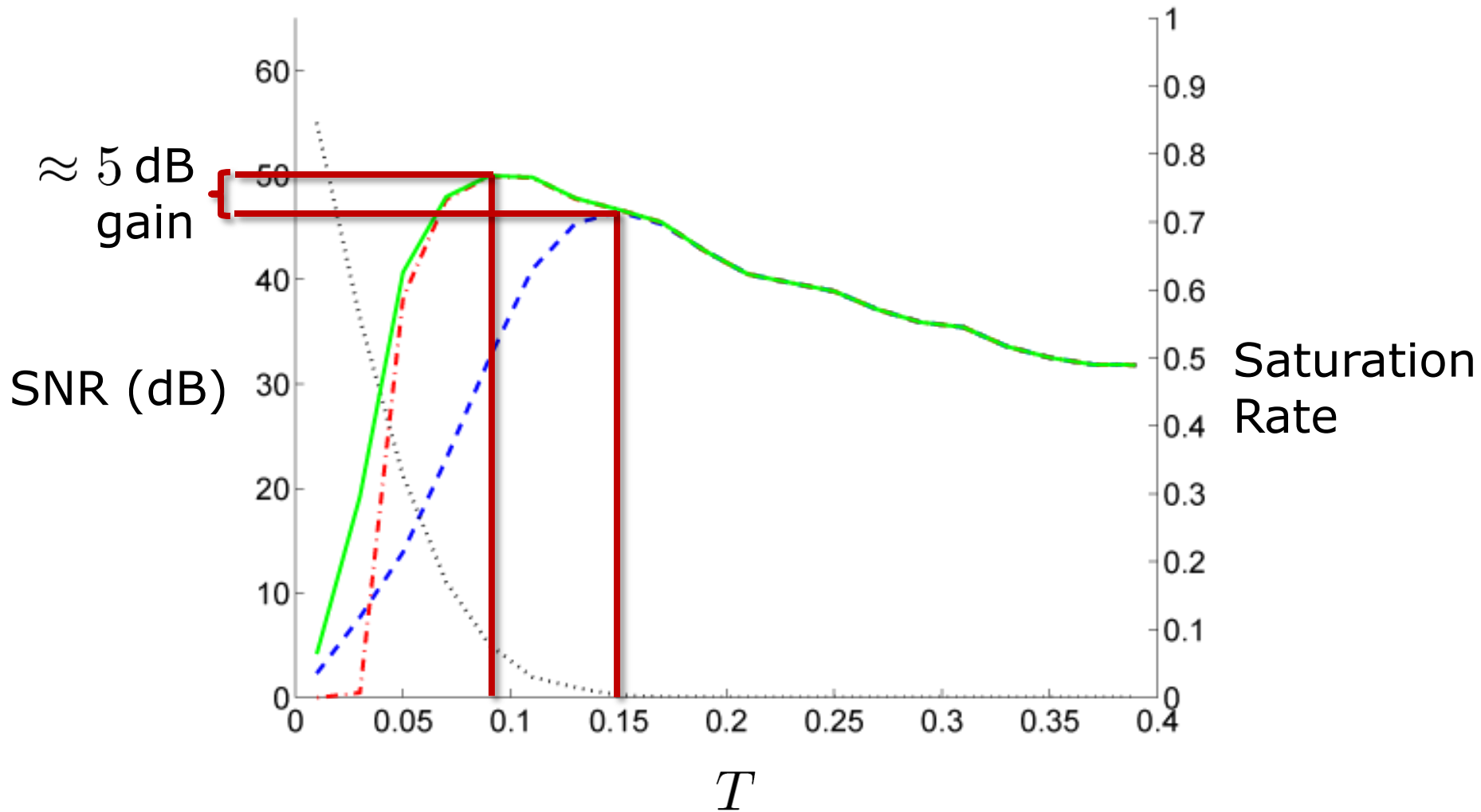
- The fact that $[\Phi \ I]$ satisfies the RIP implies that if we take D extra measurements, then we can delete $O(D)$ arbitrary rows of Φ and retain the RIP

Rejection In Practice



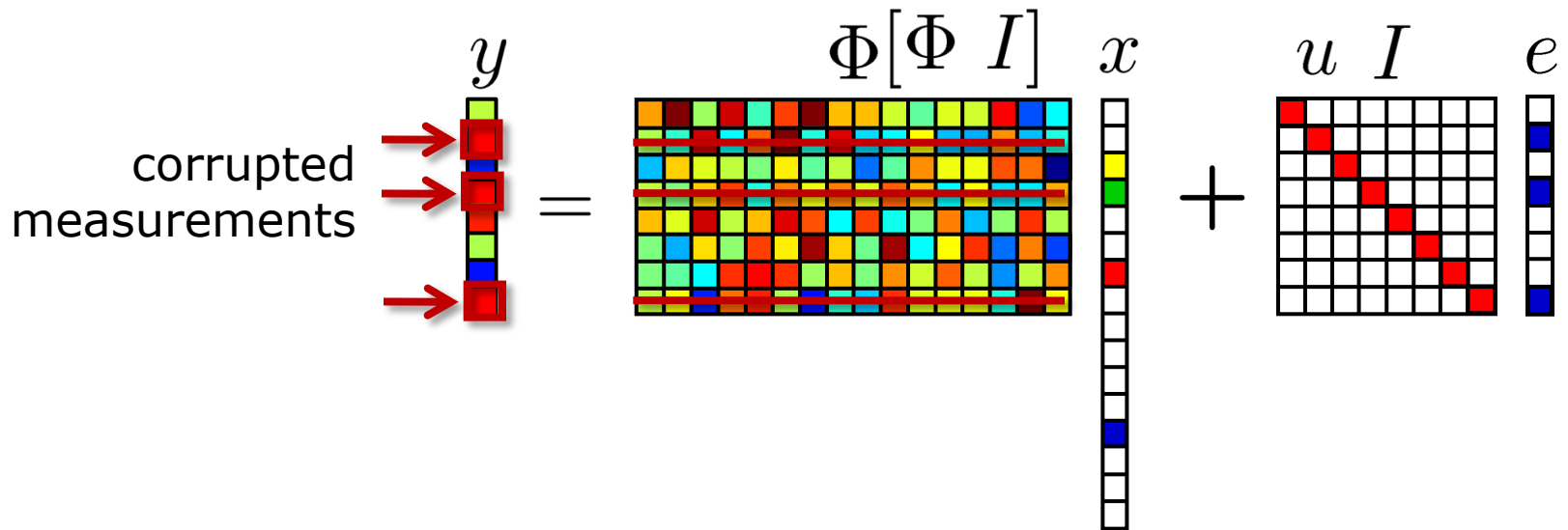
$$\text{SNR} = 10 \log_{10} \left(\frac{\|x\|_2^2}{\|\hat{x} - x\|_2^2} \right)$$

Benefits of Saturation?



Recovery in Structured Noise

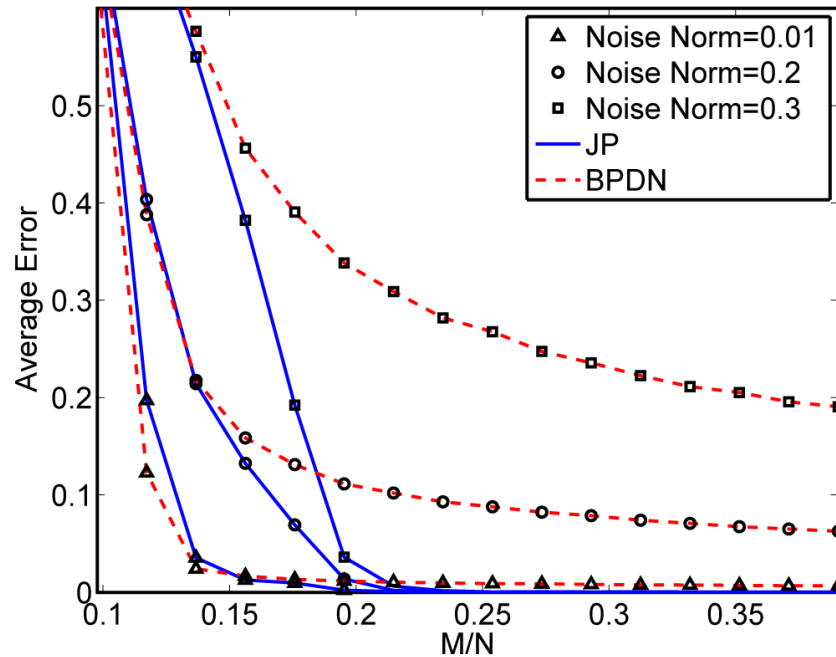
What about structured measurement noise?



Justice Pursuit

- Since $[\Phi \ I]$ satisfies the RIP, we can apply standard sparse recovery algorithms to recover u

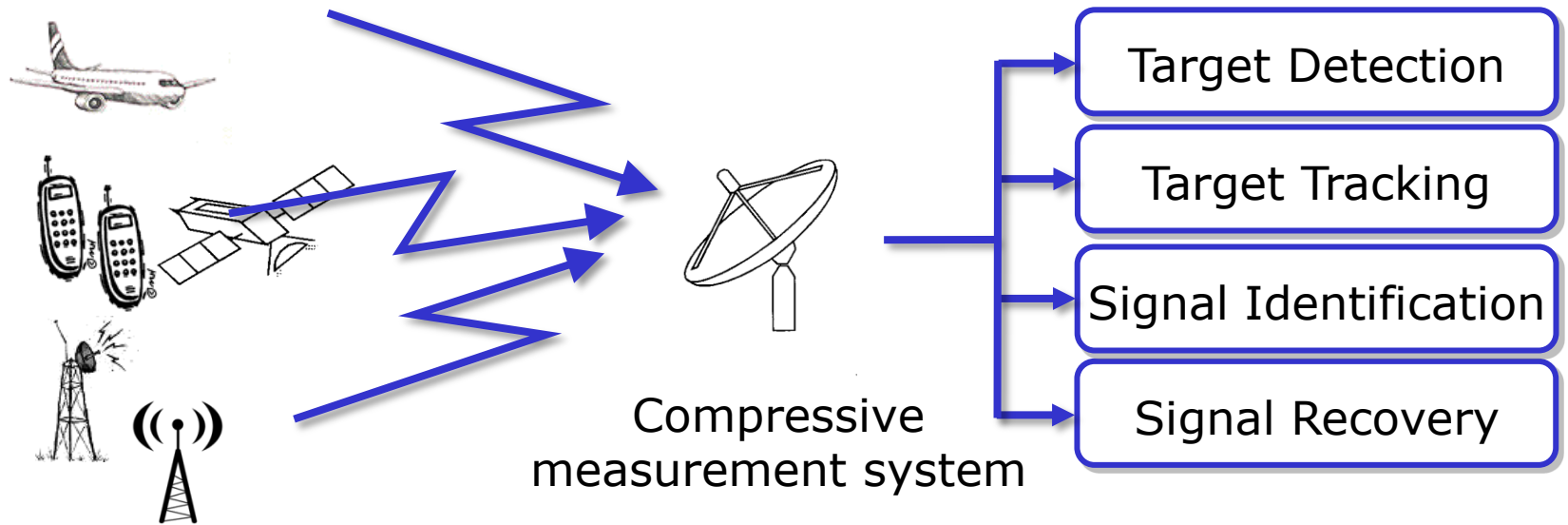
Fixed $\kappa = 10$



Compressive Signal Processing

Compressive Signal Processing

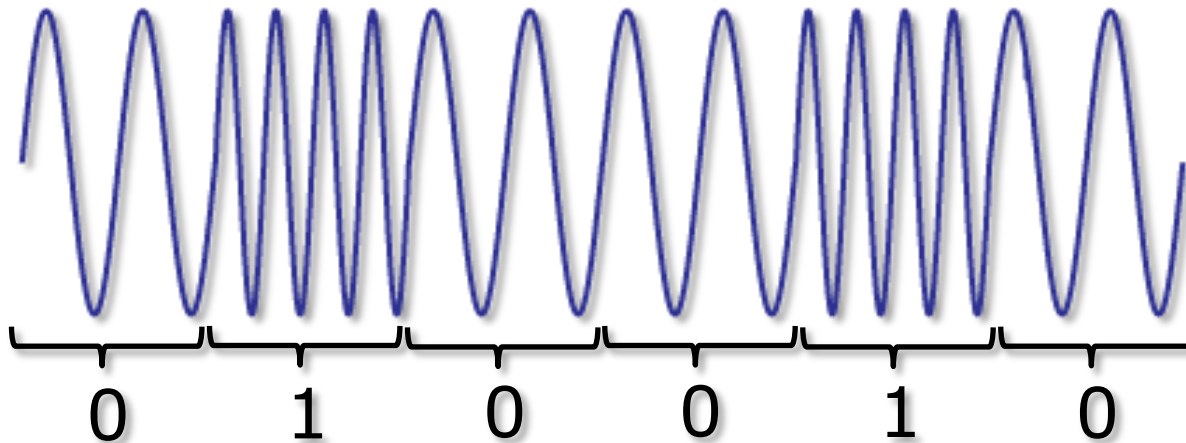
Random measurements are *information scalable*



When and how can we directly solve signal processing problems directly from compressive measurements?

Example: FM Signals

- Can we directly recover a *baseband voice signal* without recovering the modulated waveform?
- Suppose we have compressive measurements of a digital communication signal (FSK modulated)

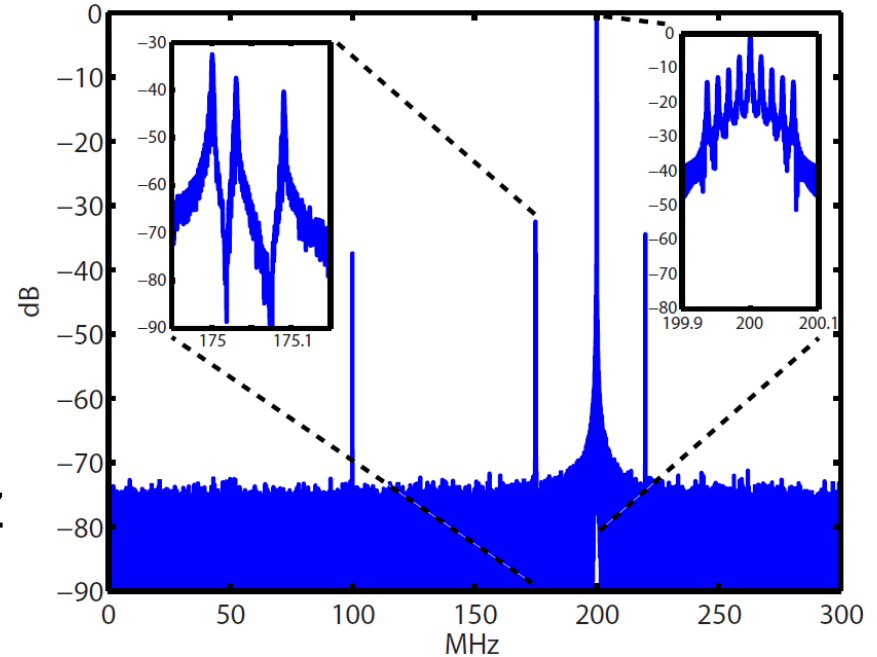


- Can we directly recover the encoded *bitstream* without first recovering the measured waveform?

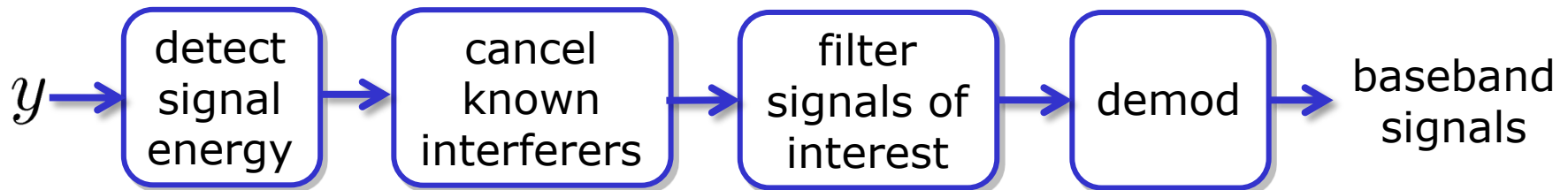
Compressive Radio Receivers

Example Scenario

- 300 MHz bandwidth
- 5 FM signals (12 kHz)
- TV station interference
- Acquire compressive measurements at 30 MHz (20 x undersampled)



We must simultaneously solve several problems

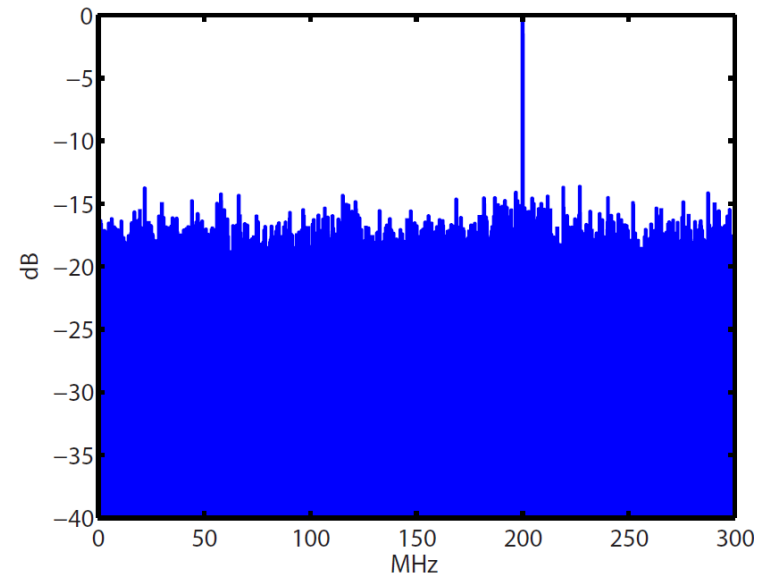
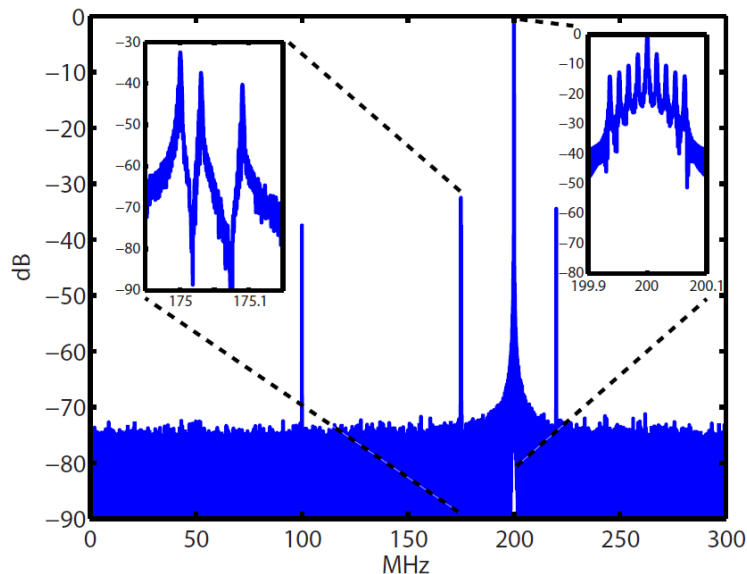


Energy Detection

We need to identify where in frequency the important signals are located

Correlate measurements with projected tones

$$\hat{F}(k) = |\langle \Phi \cos(2\pi f_k t), y \rangle|$$

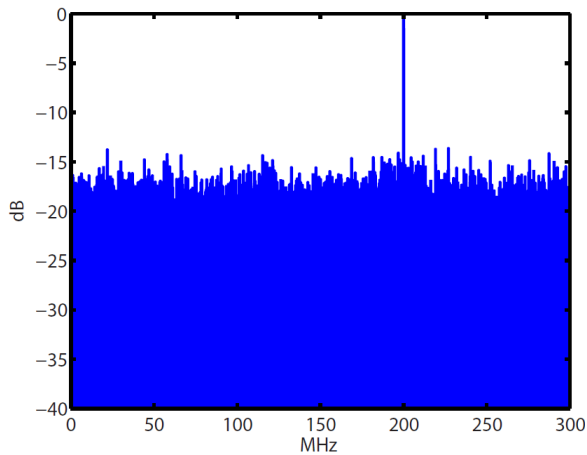


Filtering

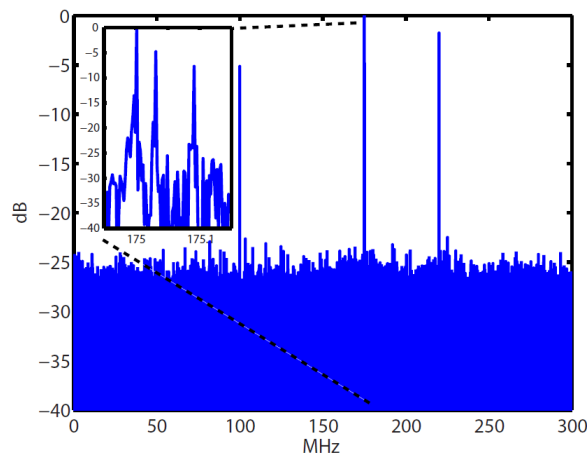
If we have multiple signals, must be able to filter to isolate and cancel interference

$$P = I - \Phi S(\Phi S)^\dagger$$

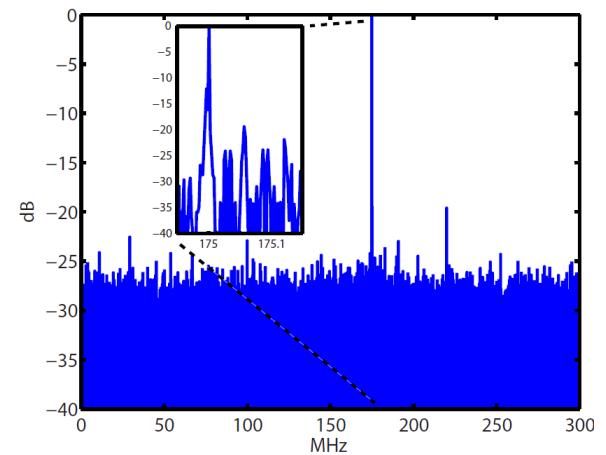
S : Discrete prolate spheroidal sequences



original



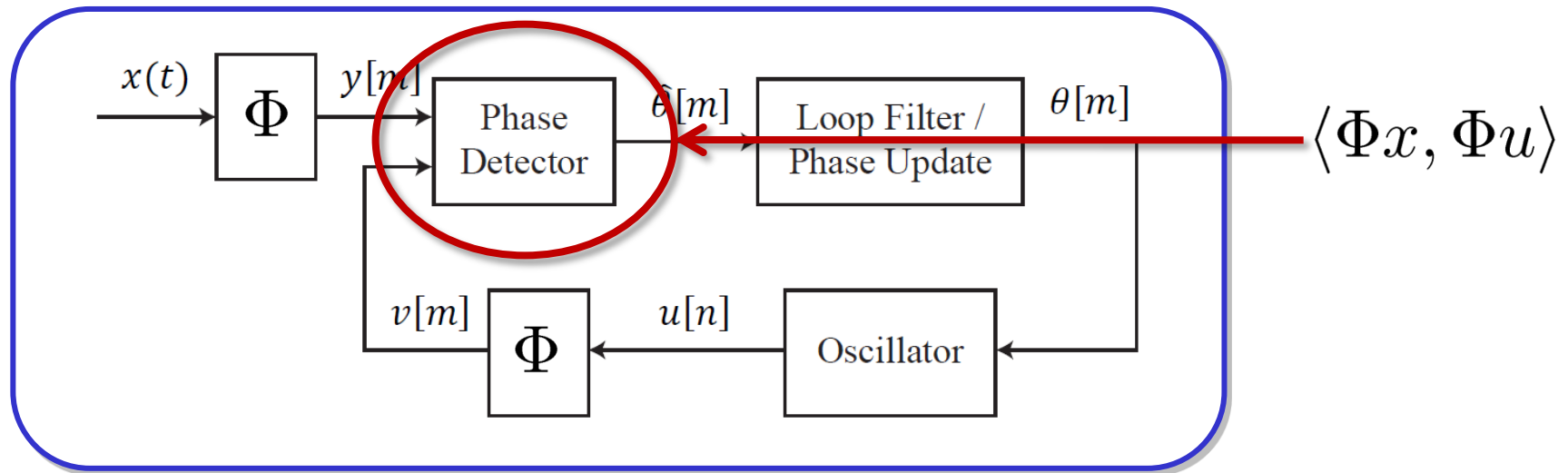
after interference
cancellation



after isolation
filtering

Unsynchronized Demodulation

We can use a phase-locked-loop (PLL) to track deviations in frequency by directly operating on compressive measurements



We can directly demodulate signals from compressive measurements *without recovery*

Summary

- **Compressive signal processing**
 - integrates sensing, compression, processing
 - exploits signal sparsity/compressibility
 - enables new sensing modalities, architectures, systems
 - exploits randomness at many levels
- Why CSP works: preserves information in signals with concise geometric structure
sparse signals | manifolds | low-dimensional models
- **Information scalability** for compressive inference
 - compressive measurements \sim sufficient statistics
 - much less computation required than for recovery

More Information

<http://dsp.rice.edu/~md>

md@rice.edu

Acknowledgements

- Rich Baraniuk
- Ron DeVore
- Piotr Indyk
- Mark Embree
- Kevin Kelly
- Petros Boufounos
- Marco Duarte
- John Treichler
- Mike Wakin
- Chinmay Hegde
- Jason Laska
- Stephen Schnelle

“I not only use all the brains
I have, but all I can borrow.”
-Woodrow Wilson

